# CS1 1212 - Spring 2022 <br> Data Structures \& <br> Introduction to Algorithms 

Data Structures<br>Trees: Tree Traversals

Ibrahim Albluwi



## Tree Data Structures

Definitions and properties
Basic operations
Balanced binary search trees

- Tree traversals


## Printing the Tree (in order)



Assuming the tree is a binary search tree, how can we traverse it in order?

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes
(recursion $\vee)$


Assuming the tree is a binary search tree, how can we traverse it in order?

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at 6
do not print 6 yet!
print left subtree first.

Console
at (6) left - current - right

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at 2
do not print 2 yet!
print left subtree first.

```
at (2) left - current - right
at (6) left - current - right
    stack frames
```


## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at 1
do not print 1 yet!
print left subtree first.

```
at (1) left - current - right
at (2) left - current - right
at (6) left - current - right
    stack frames
```


## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at NULL nothing to be done!

```
at (NULL) do nothing!
    at (1) left - current - right
    at (2) left - current - right
    at (6) left - current - right
    stack frames
```

Console

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at 1 ready to print 1

```
at (1) left - current - right
at (2) left - current - right
at (6) left - current - right

Console

\section*{Printing the Tree (in order)}

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at NULL nothing to be done
```

at (NULL) do nothing!
at (1) left - current - right
at (2) left - current - right
at (6) left - current - right
stack frames

```

Console

\section*{Printing the Tree (in order)}

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

done with tree rooted at 1
at (1) left - current - right
at (2) left - current - right
at (6) left - current - right
stack frames

Console

\section*{Printing the Tree (in order)}

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print the tree rooted at 2 ready to print 2
right subtree still needs to be printed.

Console
```

at (2) left - current - right
at (6) left - current - right

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
print tree rooted at 4
do not print 4 yet!
print left subtree first.
```

at (4) left - current - right
at (2) left - current - right
at (6) left - current - right
stack frames

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print tree rooted at 3 ready to print 3 after going left
at (3) left - current - right
at (4) left - current - right
at (2) left - current - right
at (6) left - current - right

Console
123

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
print tree rooted at 4
ready to print 4
right subtree still needs to be printed.
```

at (4) left - current - right
at (2) left - current - right
at (6) left - current - right
stack frames

Console

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print tree rooted at 5 ready to print 5 after going left
at (5) left - current - right
at (4) left - current - right
at (2) left - current - right
at (6) left - current - right
Console
12345

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
at (4) left - current - right
at (2) left - current - right
at (6) left - current - right
    stack frames
```

Console
12345

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
at (2) left - current - right
at (6) left - current - right
    stack frames
```

Console
12345

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


Console
at (6) left - current - right

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
print tree rooted at 8
do not print }8\mathrm{ yet!
print left subtree first.
```

```
at (8) left - current - right
at (6) left - current - right
```

Console

$$
123456
$$

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print tree rooted at 7 ready to print 7 after going left

```
at (7) left - current - right
at (8) left - current - right
at (6) left - current - right
```

Console

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
print tree rooted at 8
ready to print }
right subtree still needs to be printed
```

Console

```
12345678
```

```
12345678
```

```
at (8) left - current - right
at (6) left - current - right
```


## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

print tree rooted at 9 ready to print 9 after going left
at (9) left - current - right
at (8) left - current - right
at (6) left - current - right

Console

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes

done with tree rooted at 8

Console

```
123456789
```

```
123456789
```

```
at (8) left - current - right
at (6) left - current - right
```


## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


Console
at (6) left - current - right

## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
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Console

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Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
template <class T>
void BST<T>::print_in_order() const {
    print_in_order(root);
}
template <class T>
void BST<T>::print_in_order(Node<T>* node) const {
    if (node == nullptr) return;
    print_in_order(node->left);
    cout << node->val << " ";
    print_in_order(node->right);
}
```


## Printing the Tree (in order)

Idea. Print all of the left subtree and then print the current node and then print all of the right subtree.
I.e. Print the smaller nodes then print the current node, then print the larger nodes


```
template <class T>
void BST<T>::print_in_order() const {
    print_in_order(root);
}
```

template <class T>
void BST<T>::print_in_order(Node<T>* node) const \{
if (node == nullptr) return;
_ private helper
(recursive) function
print_in_order(node->left);
cout << node->val << " ";
print_in_order(node->right);
\}

## Clearing the Tree



How can we traverse the tree and delete every node?

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


How can we traverse the tree and delete every node?

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree Clear the right subtree, then delete the current node


To clear the tree rooted at 6, 2 and 8 have to be cleared first

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


To clear the tree rooted at 2, 1 and 4 have to be cleared first

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


1 can be cleared!

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


To clear the tree rooted at 4, 3 and 5 have to be cleared first

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


1 and 5 can be cleared!

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


4 can be cleared!

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
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Clear the right subtree, then delete the current node


## Clearing the Tree

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Clear the right subtree, then delete the current node


To clear the tree rooted at 8,
7 and 9 have to be cleared first

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


7 and 9 can be cleared!

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree

Clear the right subtree, then delete the current node


6 can be cleared!

## Clearing the Tree

Idea. Since the children of a node are accessible only through the node, do not delete the node until its children have been deleted.
I.e. Clear the left subtree Clear the right subtree, then delete the current node


```
template <class T>
void BST<T>::clear() {
    clear(root);
    root = nullptr;
}
template <class T>
void BST<T>::clear(Node<T>* node) {
    if (node == nullptr) return;
    clear(node->left);
    clear(node->right);
    delete node;
}
```


## Copying the Tree

BST Copy Constructor. Given another BST named other, insert every node in other into the current tree, such that the current tree becomes exactly like other.

Procedure. Insert the current node value Copy the left subtree Copy the right subtree


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Procedure. Insert the current node value Copy the left subtree Copy the right subtree


```
template <class T>
BST<T>::BST(const BST<T>& other) {
    root = nullptr;
    copy_from(other.root);
}
template <class T>
void BST<T>::copy_from(Node<T>* node) {
    if (node == nullptr) return;
    insert(node->val);
    copy_from(node->left);
    copy_from(node->right);
}
```

! Running Time. $n$ nodes are inserted.
In general: $\quad O(n \times h e i g h t)$ Balanced Trees: $O(n \log n)$ Worst Case: $\quad O\left(n^{2}\right)$
! Note. This code creates an exact copy of the tree assuming the insert function does not rebalance with rotations!

## Depth-First Traversals

Problem. Given a binary tree, visit every node in the tree and perform some operation (e.g. delete the node, print the node, etc.)

Solution. Start at the root and use recursion to traverse the left and right subtrees.

Pre-order Traversal. Perform the operation before traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    do_something();
    preOrder(node->left);
    preOrder(node->right);
```

In-order Traversal. Perform the operation after traversing the left subtree and before traversing the right subtree.

```
void inOrder(Node<T>* node) {
    if (node == nullptr) return;
    inOrder(node->left);
    do_something();
    inOrder(node->right);
}
```

Post-order Traversal. Perform the operation after traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    preOrder(node->left);
    preOrder(node->right);
    do_something();
}
```


## Depth-First Traversals

Problem. Given a binary tree, visit every node in the tree and perform some operation (e.g. delete the node, print the node, etc.)

Solution. Start at the root and use recursion to traverse the left and right subtrees.

Example applications. Copying a BST and computing node depths.

Pre-order Traversal. Perform the operation before traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    do_something();
    preOrder(node->left);
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}
```

In-order Traversal. Perform the operation after traversing the left subtree and before traversing the right subtree.

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void inOrder(Node<T>* node) {
    if (node == nullptr) return;
    inOrder(node->left);
    do_something();
    inOrder(node->right);
}
```

Post-order Traversal. Perform the operation after traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    preOrder(node->left);
    preOrder(node->right);
    do_something();
}
```


## Depth-First Traversals

Problem. Given a binary tree, visit every node in the tree and perform some operation (e.g. delete the node, print the node, etc.)

Solution. Start at the root and use recursion to traverse the left and right subtrees.

> Example application. Printing a BST in order

Pre-order Traversal. Perform the operation before traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    do_something();
    preOrder(node-> left);
    preOrder(node->right);
}
```

In-order Traversal. Perform the operation after traversing the left subtree and before traversing the right subtree.

```
void inOrder(Node<T>* node) {
    if (node == nullptr) return;
    inOrder(node->left);
    do_something();
    inOrder(node->right);
}
```

Post-order Traversal. Perform the operation after traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    preOrder(node->left);
    preOrder(node->right);
    do_something();
}
```


## Depth-First Traversals

Problem. Given a binary tree, visit every node in the tree and perform some operation (e.g. delete the node, print the node, etc.)

Solution. Start at the root and use recursion to traverse the left and right subtrees.

Example applications. Clearing a tree and computing node heights.

Pre-order Traversal. Perform the operation before traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    do_something();
    preOrder(node->left);
    preOrder(node->right);
}
```

In-order Traversal. Perform the operation after traversing the left subtree and before traversing the right subtree.

```
void inOrder(Node<T>* node) {
    if (node == nullptr) return;
    inOrder(node->left);
    do_something();
    inOrder(node->right);
}
```

Post-order Traversal. Perform the operation after traversing the left and right subtrees.

```
void preOrder(Node<T>* node) {
    if (node == nullptr) return;
    preOrder(node->left);
    preOrder(node->right);
    do_something();
}
```


## Printing the Tree (level-by-level)



How can we traverse the tree level-by-level?

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Queue

Console

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Remove $\mathbf{A}$ and add its children


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Remove B and add its children

## Queue



Console


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Remove $\mathbf{C}$ and add its children

## Queue



## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Remove $\mathbf{D}$ and add its children

## Queue



Console
A B C D

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Remove E and add its children


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Console
ABCDEFG

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


Console
ABCDEFGH

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Queue


Console
ABCDEFGHIJ

## Printing the Tree (level-by-level)

Idea. Maintain a queue of the nodes yet to be visited.

Repeat until the queue is empty: Remove a node and add its children.


## Printing the Tree (level-by-level)

```
template <class T>
void BST<T>::level_order() const {
    if (is_empty()) return;
    Queue<Node<T>* > queue;
    queue.enqueue(root);
    while (!queue.is_empty()) {
        Node<T>* node = queue.dequeue();
        cout << node->val << " ";
        if (node->left != nullptr) queue.enqueue(node->left);
        if (node->right != nullptr) queue.enqueue(node->right);
    }
}
```


## Printing the Tree (level-by-level)

```
template <class T>
void BST<T>::level_order() const {
    if (is_empty()) return;
    Queue<Node<T>*> queue;
    queue.enqueue(root);
    while (!queue.is_empty()) {
        Node<T>* node = queue.dequeue();
        cout << node->val << " ";
        if (node->left != nullptr) queue.enqueue(node->left);
        if (node->right != nullptr) queue.enqueue(node->right);
    }
}
```


## Printing the Tree (level-by-level)

```
template <class T>
void BST<T>::level_order() const {
    if (is_empty()) return;
    Queue<Node<T>* > queue;
    queue.enqueue(root);
    while (!queue.is_empty()) {
        Node<T>* node = queue.dequeue();
        cout << node->val << " ";
        if (node->left != nullptr) queue.enqueue(node->left);
        if (node->right != nullptr) queue.enqueue(node->right);
    }
}
```


## Printing the Tree (level-by-level)

```
template <class T>
void BST<T>::level_order() const {
    if (is_empty()) return;
    Queue<Node<T>*> queue;
    queue.enqueue(root);
    while (!queue.is_empty()) {
        Node<T>* node = queue.dequeue();
        cout << node->val << " ";
        if (node-> left != nullptr) queue.enqueue(node-> left);
        if (node->right != nullptr) queue.enqueue(node->right);
    }
}
```


## Printing the Tree (level-by-level)

```
template <class T>
void BST<T>::level_order() const {
    if (is_empty()) return;
    Queue<Node<T>*> queue;
    queue.enqueue(root);
    while (!queue.is_empty()) {
        Node<T>* node = queue.dequeue();
        cout << node->val << " ";
        if (node-> left != nullptr) queue.enqueue(node-> left);
        if (node->right != nullptr) queue.enqueue(node->right);
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}
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Terminology. Breadth-First Traversal (BFT) $=$ Level-Order Traversal

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}
```

Terminology. Breadth-First Traversal (BFT) = Level-Order Traversal Right-to-left BFT. Enqueue the right child before the left child.

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}
```

Terminology. Breadth-First Traversal (BFT) = Level-Order Traversal
Right-to-left BFT. Enqueue the right child before the left child.

Note. While in this code we print each dequeued node, the BFT is a general-purpose traversal that can be used to go through all the nodes in the tree and perform some operation (e.g. printing the node value, computing and storing the node depth, checking if the node is a leaf, etc.)

## What does the following function do?

```
template <class T>
void BST<T>::mystery() const {
    if (is_empty()) return;
    Stack<Node<T>*> stack;
    stack.push(root);
    while (!stack.is_empty()) {
        Node<T>* node = stack.pop();
        cout << node->val << " ";
        if (node->right != nullptr)
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stack

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            stack.push(node->right);
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Console

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    }
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stack

Console

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if (node->right != nullptr)
stack.push(node->right);
if (node->left != nullptr)
stack.push(node-> left);
}
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```

stack
Console

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if (is_empty()) return;
Stack<Node<T>* > stack;
stack.push(root);
while (!stack.is_empty()) {
Node<T>* node = stack.pop();
cout << node->val << " ";
- if (node->right != nullptr)
stack.push(node->right);
if (node->left != nullptr)
stack.push(node->left);
}
}

- if (node->right != nullptr) stack. push(node $->$ right);



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        Node<T>* node = stack.pop();
        cout << node->val << " ";
    }
}
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            stack.push(node->right);
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            stack.push(node->left);
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stack
Console

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```


stack

Console

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stack

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void BST<T>::mystery() const {
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        Node<T>* node = stack.pop();
        cout << node->val << " ";
        if (node->right != nullptr)
            stack.push(node->right);
        if (node->left != nullptr)
            stack.push(node->left);
    }
}
Pre-order Traversal!
```

Console

```
6 2 14 3 5 8 7 9
```


## More Practice Exercises

1. Store in every node its height.
2. Store in every node its depth.
3. Count the number of nodes in the tree or count the number of leafs in the tree.
4. Store in every node the number of nodes in the subtree rooted at that node.
5. Find the maximum in a general binary tree (not a BST)
6. Print the tree in reverse level-order (from the right-most node in the last level to the root).
7. Find the median in a BST (in $\mathrm{O}(n)$ )
8. Find the median in a balanced BST (in $\mathrm{O}(\log n)$ ) assuming exercise 4 is solved.
9. Remove all the leafs from the tree.
10. Count all the nodes in the last level.
... and many more!

## Back to the Set ADT

Problem. Design a data structure to support the following operations:

- insert(val) // add val to the set if it is not already in the set.
- remove(val) // remove val from the set of items.
- contains(val) // check if val belongs to the set.

Candidate implementations.

|  | insert(val) | remove(val) | contains (val) |
| :---: | :---: | :---: | :---: |
| Unordered DLL | O(n) | O(n) | O(n) |
| Unordered SLL | O(n) | O(n) | O(n) |
| Ordered DLL | O(n) | 0 (n) | O(n) |
| Ordered SLL | 0 ( n ) | 0 ( n ) | O(n) |
| Unordered Array | O(n) | O(n) | O(n) |
| Ordered Array | O(n) | 0 (n) | $O(\log n)$ |
| Balanced BST | $0(\log n)$ | $0(\log n)$ | $0(\log n)$ |

## Another ADT: A Map (or Dictionary)

Problem. Design a data structure to support the following operations:

- insert(key, val) // insert a new key-value pair or reset // the current value of they key
- remove(key) // remove the key and its corresponding value
- get(key) // return the value corresponding to the key


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- insert(key, val) // insert a new key-value pair or reset // the current value of they key
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Example Applications.

- A mapping between words and their meanings (key and val are string)
- A mapping between usernames and passwords (key and val are string)
- A mapping between IDs and GPAs (key is string and val is float)
- A mapping between years and number of new borns (key and val are int)


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- A mapping between years and number of new borns (key and val are int)

Solution. Use a balanced BST. Modify the Node class to have a key and a value.

- insert(key, val) // search based on key. If key is found, change the // current val, else insert a new node $\longrightarrow 0(l o g n)$
- remove(key) // same as remove in the set ADT $\longrightarrow 0(\log n)$
- get(key) // search based on key $\longrightarrow 0(\log n)$

