CS11212 - **Spring** 2022

Data Structures & Introduction to **Algorithms**

Data Structures

Trees: Definitions and Properties

Ibrahim Albluwi

Problem. Design a data structure to support the following operations:

- insert(val) // add val to the set if it is not already in the set.
- remove(val) // remove val from the set of items.
- contains(val) // check if val belongs to the set.

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Candidate implementations.

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Ordered DLL	O(n)	O(n)	0(n)	Maintaining
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Ordered SLL	O(n)	O(n)	O(n)	help!
Unordered Array	O(n)	O(n)	O(n)	Bad!
Ordered Array	O(n)	O(n)	O(log n)	Great!
				are bad!

Bad!

What is going on?

- Linked lists support efficient insertion and deletion. Provided that there is a pointer to the insertion or deletion position.
- Provided that there is a pointer to the Arrays support efficient search. Provided that the array is sorted.
 - Linked lists do not support efficient search. Sorted linked lists do not support efficient binary search because it requires direct access.
 - Arrays do not support efficient insertion and removal at any position. Elements need to be shifted.

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Can we achieve the good of the two?

Efficient search (as in binary search)

Efficient insertion / removal (given a pointer to the insertion / removal position)

Bad!

$0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 10 \rightarrow 11 \rightarrow 12 \rightarrow 13 \rightarrow 14$

Can we perform *binary search* on a sorted linked list?



O(n) operations to get to the middle element!



Idea: Maintain a pointer to the middle element.

The middle element is now accessible in O(1).



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Still O(n) to get to the middle of the left half or the middle of the right half!



Idea (again): Maintain a pointer to the middle elements!



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Idea (*again*, *again*): Maintain a pointer to the middle elements!



Idea (*again*, *again*): Maintain a pointer to the middle elements!



A Binary Search Tree!

Proposed Implementation



A Binary Search Tree!

We need to build a <u>linked</u> *tree structure*.

Proposed Implementation



My family

A Binary Search Tree!

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Trees are not useful only for efficient search. They are useful for representing hierarchies and relationships in general.



Proposed Implementation



A Binary Search Tree!

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Questions.

- How do we *insert* into the tree?
- How do we *remove* from the tree?
- How do we *search* the tree?
- How do we *traverse* the tree?
- What *other operations* can we perform?
- What *properties* do trees have?
- ... ?



Tree Data Structures

Definitions and properties
 Operations on BSTs
 Balanced binary search trees
 Tree traversals

Definition. A *tree* is:

- NULL or
- a node connected to a set of *disjoint* trees.



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A node connected to two disjoint trees

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A node connected to two disjoint trees:

- A NULL tree and
- A node connected to two NULL trees

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not a tree: A node connected to two overlapping trees

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Alternative definition (from graph theory).

A *tree* is a set of nodes in which any two nodes are connected by exactly one path



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A binary tree in nature!



Binary trees?

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nodes that share the same parent are siblings
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parent, grandparent,

descendents

children, grandchildren, grand grand children, etc.

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Every tree has one root node.

The root node has no parent. The root is an ancestor for all nodes. All nodes are descendants of the root.



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Leafs are also called "external nodes".

internal nodes



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The degree of a node is the number of its children. The degree of a tree is the maximum degree of a node in the tree. A leaf has degree 0. The illustrated tree has degree 2.





The depth of a node is its level in the tree.

The root has depth 0.



The **depth** of a node is its level in the tree.

The root has depth 0.

The height of a node is the maximum number of levels below it. All leafs have height 0.



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The height of a tree is the height of the root. Also called "tree depth".



Root	=
Leafs	=
Tree height	=
Tree Degree	=
Ancestors of L	=

- Height of E =
 - Depth of E =
- Degree of E =
- Descendants of E =
 - Siblings of G =



Root = A
Leafs = I J N M F G H
Tree height = 5
Tree Degree = 3
Ancestors of L = K E B A

Height of E = 3
Depth of E = 2
Degree of E = 1
Descendants of E = K L M N
Siblings of G = F H











Why is binary search possible on such a tree?



- larger than all the nodes to its left.
- smaller than all the nodes to its right.
- a binary search tree.



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Question. Which node is always the *maximum* node in a BST?



Question. Which node is always the *maximum* node in a BST? **Answer**. The right-most node.



Question. Which node is always the *maximum* node in a BST? **Answer**. The right-most node.

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Question. Which node is always the *minimum* node in a BST? **Answer**. The left-most node.

Question. Which node is always the *median* node in a BST?



Question. Which node is always the *maximum* node in a BST? **Answer**. The right-most node.

Question. Which node is always the *minimum* node in a BST? **Answer**. The left-most node.

Question. Which node is always the *median* node in a BST? **Answer**. Impossible to tell without performing a search for the median. *stay tuned!*



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Binary Search Trees



Definition. A binary search tree (BST) is a binary tree where each node is:

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- smaller than all the nodes to its right.
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Searching a BST requires *O*(*height*) compares.

what is the height of a BST?



Definition.

- All levels are full.
- All leafs are in the last level and all internal nodes have exactly two children.

Ν = H = 0Ν = 3 H = 1Ν 7 = H = 2N = 15H = 3N = 31H = 4

CAUTION You might find different definitions online. Some texts call this tree *complete*

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of nodes (
$$N$$
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 $\text{Height} = \log_2(N+1) - 1$



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of nodes (
$$N$$
) = 2 height + 1 - 1

Height = $\log_2(N + 1) - 1$

IMPORTANT!

The height of a *perfect* binary tree is *logarithmic* in the number of nodes in the tree!



Definition.

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of nodes (*N*) =
$$2^{\text{height} + 1} - 1$$

Height = $\log_2(N + 1) - 1$

of leafs = 2 height

of leafs = $\frac{N+1}{2}$ # of internal nodes = $\frac{N-1}{2}$



Balanced Binary Tree. the heights of the two child subtrees of any node differ by at most one.

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Balance Factor of a node (bf) = height of left child – height of right child



IMPORTANT!

The height of a *balanced* binary tree is *logarithmic* in the number of nodes in the tree.



Tree Data Structures

Definitions and properties

• Basic Operations

Balanced binary search trees

Tree traversals